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**LEUVEN**

# Mining Frequent Trees in Networks

## Problem

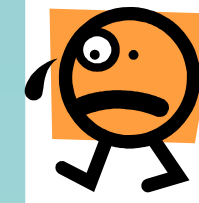
### Frequent Tree Mining:

Given a graph  $G$  and threshold  $t$ , find and list all trees that **occur** in  $G$  at least  $t$  times

#### Setting:

- Single large networks
- Subgraph isomorphism as matching operator
- Pattern: Labeled rooted trees

## Challenges



- Deciding whether a pattern is subgraph isomorphic to  $G$  itself is NP-Complete (naïve solution is  $O(n^k)$ )
- How to measure frequency?
- Main focus was to tackle computational complexity of the problem

## Proposal

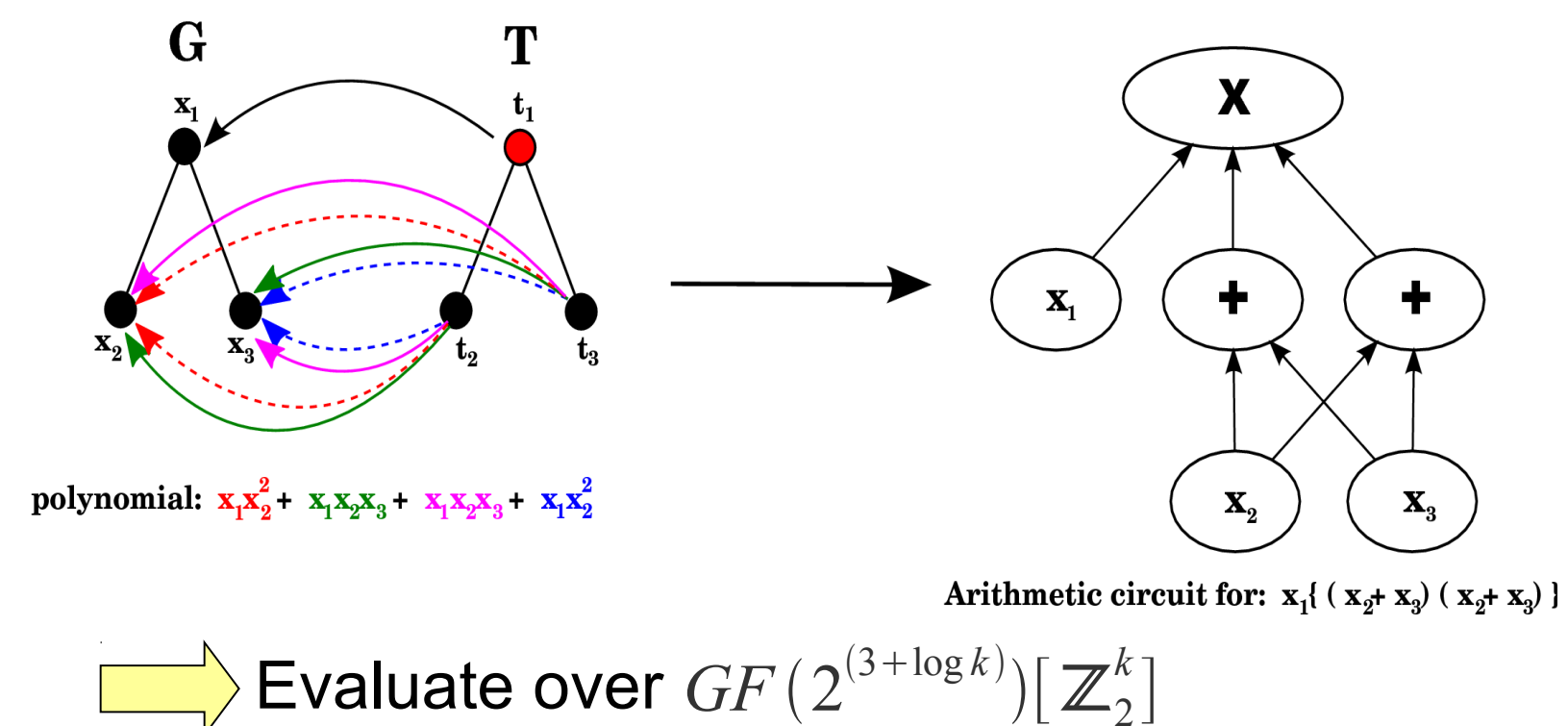
- Recent advances in parameterized complexity theory, give us (randomized) algebraic methods that for tree patterns, reduce subgraph isomorphism from  $O(n^k)$  to  $O(m \cdot k^2 \cdot 2^k \cdot \log^2 k)$

- Combined with a simple anti-monotone frequency measure, we can mine *all* tree patterns with similar complexity, in **delay**  $O(|\Sigma| \cdot m \cdot k^2 \cdot 2^k \cdot \log^2 k)$

→ **Frequency:** Images of root under isomorphism

- Accuracy can be boosted arbitrarily by repeating subgraph isomorphism multiple times
- Exponential factors are relatively small, to allow for practical applications

### Method Overview



### Delay time:

The time between any two consecutive outputs (including start and 1<sup>st</sup> output) is an unchanging delay

## Researcher

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## Promoter

Jan Ramon

## Sponsors



## Results

- 1) What pattern/networks size can we handle within reasonable time?  
→ up to trees size 15 in  $10^4$  and size 10 in  $10^6$  networks
- 2) How does our strategy compare to state of the art methods?  
→ outperforms VF2 even for moderate size problems
- 3) Does implementation really scale as well as given theoretical bounds?  
→ approximately, yes
- 4) What is the influence of different parameters and optimizations?  
→ See right

