

# A Belgian view on lattice rules

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Linz, Austria, October 14–18, 2013



## Atomium – Brussels

built in 1958

height  $\approx 103\text{m}$

figure = 2€ coin  
 $5 \cdot 10^6$  in circulation

Body centered cubic lattice

# Introduction

Given is an integral

$$I[f] := \int_{\Omega} w(\mathbf{x})f(\mathbf{x}) \, d\mathbf{x}$$

where  $\Omega \subseteq \mathbb{R}^s$  and  $w(\mathbf{x}) \geq 0, \forall \mathbf{x} \in \mathbb{R}^s$ .

Search an approximation for  $I[f]$

$$I[f] \simeq Q[f] := \sum_{j=1}^n w_j f(\mathbf{y}^{(j)})$$

with  $w_j \in \mathbb{R}$  and  $\mathbf{y}^{(j)} \in \mathbb{R}^s$ .

Webster:

**quadrature**: the process of finding a square equal in area to a given area.

**cubature**: the determination of cubic contents.

If  $s = 1$  then  $Q$  is called a **quadrature formula**.

If  $s \geq 2$  then  $Q$  is called a **cubature formula**.

$$Q[f] := \sum_{j=1}^n w_j f(\mathbf{y}^{(j)})$$

Cubature/quadrature formulas are **basic integration rules**

→ choose points  $\mathbf{y}^{(j)}$  and weights  $w_j$  independent of integrand  $f$ .

It is difficult (time consuming) to construct basic integration rules, but the result is usually hard coded in programs or tables.

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**Restriction to unit cube:** given is

$$I[f] = \int_0^1 \cdots \int_0^1 f(x_1, \dots, x_s) dx_1 \cdots dx_s = \int_{[0,1]^s} f(\mathbf{x}) d\mathbf{x}$$

Taxonomy: two major classes

- 1 polynomial based methods  
incl. methods exact for algebraic or trigonometric polynomials
- 2 number theoretic methods  
incl. Monte Carlo and quasi-Monte Carlo methods

As in zoology, some species are difficult to classify.

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For example

### Definition

An  $s$ -dimensional lattice rule is a cubature formula which can be expressed in the form

$$Q[f] = \frac{1}{d_1 d_2 \dots d_t} \sum_{j_1=1}^{d_1} \sum_{j_2=1}^{d_2} \dots \sum_{j_t=1}^{d_t} f \left( \left\{ \frac{j_1 \mathbf{z}_1}{d_1} + \frac{j_2 \mathbf{z}_2}{d_2} + \dots + \frac{j_t \mathbf{z}_t}{d_t} \right\} \right),$$

where  $d_i \in \mathbb{N}_0$  and  $\mathbf{z}_i \in \mathbb{Z}^s$  for all  $i$ .

Alternative formulation:

### Definition

A multiple integration lattice  $\Lambda$  is a subset of  $\mathbb{R}^s$  which is discrete and closed under addition and subtraction and which contains  $\mathbb{Z}^s$  as a subset.

### Definition

A **lattice rule** is a cubature formula where the  $n$  points are the points of a **multiple integration lattice**  $\Lambda$  that lie in  $[0, 1)^s$  and the weights are all equal to  $1/n$ .

$$n = n(Q) = \#\{\Lambda \cap [0, 1)^s\} .$$



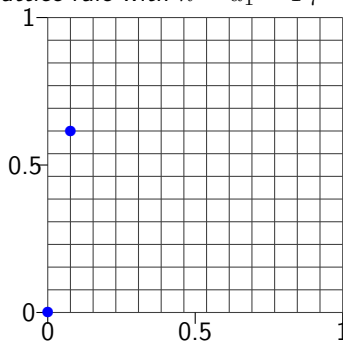
# Example

The Fibonacci lattice with  $n = F_j$  and  $\mathbf{z} = (1, F_{j-1})$

has points  $\mathbf{x}^{(j)} = \left( \frac{j}{F_j}, \frac{jF_{j-1}}{F_j} \right)$

$\Rightarrow$  lattice rule  $Q[f] = \frac{1}{n} \sum_{j=0}^{n-1} f \left( \left\{ \frac{(j, jF_{j-1})}{n} \right\} \right)$

Example: the lattice rule with  $n = d_1 = F_7 = 13$  and  $\mathbf{z}_1 = (1, 8)$



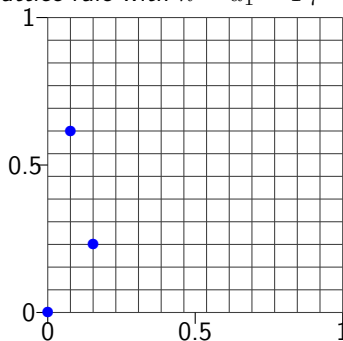
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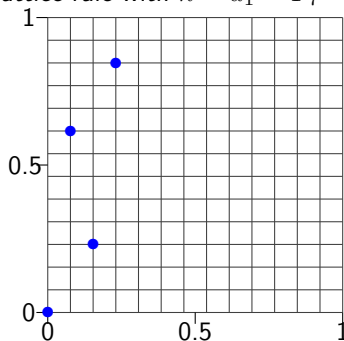
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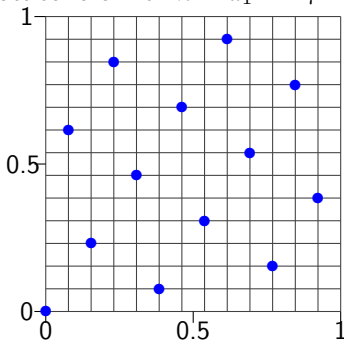
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# Polynomials

Let  $\alpha = (\alpha_1, \alpha_2, \dots, \alpha_s) \in \mathbb{Z}^s$  and  $|\alpha| := \sum_{j=1}^s |\alpha_j|$ .

algebraic polynomial

$$p(\mathbf{x}) = \sum a_\alpha \mathbf{x}^\alpha = \sum a_\alpha \prod_{j=1}^s x_j^{\alpha_j}, \quad \text{with } \alpha_j \geq 0$$

trigonometric polynomial

$$t(\mathbf{x}) = \sum a_\alpha e^{2\pi i \alpha \cdot \mathbf{x}} = \sum a_\alpha \prod_{j=1}^s e^{2\pi i x_j \alpha_j}$$

The **degree** of a polynomial =  $\max_{a_\alpha \neq 0} |\alpha|$ .

$\mathbb{P}_d^s$  = all algebraic polynomials in  $s$  variables of degree at most  $d$ .

$\mathbb{T}_d^s$  = all trigonometric polynomials in  $s$  variables of degree at most  $d$ .

# Quality criteria?

## Definition

A cubature formula  $Q$  for an integral  $I$  has algebraic (trigonometric) degree  $d$  if it is exact for all polynomials of algebraic (trigonometric) degree at most  $d$ .

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How many points are needed in a cubature formula to obtain a specified degree of precision?

The dimensions of the vector spaces of polynomials are:

$$\dim \mathbb{P}_d^s = \binom{s+d}{d}$$

$$\dim \mathbb{T}_d^s = \sum_{j=0}^s \binom{s}{j} \binom{d}{j} 2^j.$$

We will use the symbol  $\mathbb{V}_d^s$  to refer to one of the vector spaces  $\mathbb{P}_d^s$  or  $\mathbb{T}_d^s$ .



## Theorem

*If a cubature formula is exact for all polynomials of  $\mathbb{V}_{2k}^s$ , then the number of points  $n \geq \dim \mathbb{V}_k^s$ .*

Algebraic degree: For  $s = 2$  (Radon, 1948); general  $s$  (Stroud, 1960)  
Trigonometric degree: (Mysovskikh, 1987)



J. Radon



A. Stroud



И.П. Мысовских

## Theorem

*If a cubature formula is exact for all polynomials of degree  $d > 0$  and has only real points and weights, then it has at least  $\dim \mathbb{V}_k^s$  positive weights,  $k = \lfloor \frac{d}{2} \rfloor$ .*

Algebraic degree: (Mysovskikh, 1981)

Trigonometric degree: (C. 1997)

⇒ minimal formulas have only positive weights.

## Corollary

*If a cubature formula of trigonometric degree  $2k$  has  $n = \dim \mathbb{T}_k^s$  points, then all weights are equal.*

This is a reason to restrict searches to

$$Q[f] = \frac{1}{n} \sum_{j=1}^n f(\mathbf{x}_j).$$

# Improved bound for odd degrees

For algebraic degree, the improved lower bound for odd degrees takes into account the symmetry of the integration region.

E.g., centrally symmetric regions such as a cube  
→ (Möller, 1973)



H.M. Möller

Result for trigonometric degree is very similar.

# Improved bound for odd degrees

$G_k$  := span of trigonometric monomials of degree  $\leq k$   
with the same parity as  $k$ .

Theorem ((Noskov, 1985), (Mysovskikh, 1987))

*The number of points  $n$  of a cubature formula for the integral over  $[0, 1)^s$  which is exact for all trigonometric polynomials of degree at most  $d = 2k + 1$  satisfies*

$$n \geq 2 \dim G_k.$$

## Definition

A cubature formula is called **shift symmetric** if it is invariant w.r.t. the group of transformations

$$\left\{ \mathbf{x} \mapsto \mathbf{x}, \mathbf{x} \mapsto \left\{ \mathbf{x} + \left( \frac{1}{2}, \dots, \frac{1}{2} \right) \right\} \right\}$$

(This is the 'central symmetry' for the trig. case.)

## Theorem (Beckers & C., 1993)

*If a shift symmetric cubature formula of degree  $2k + 1$  has  $n = 2 \dim G_k$  points, then all weights are equal.*

## Conjecture (C., 1997)

*Any cubature formula that attains the lower bound is shift symmetric.*

This became a Theorem (Osipov, 2001).

# Known minimal formulas for trigonometric degree

- for all  $s$ 
  - degree 1
  - degree 2 (Noskov, 1988)
  - degree 3 (Noskov, 1988)
- for  $s = 2$ 
  - all even degrees (Noskov, 1988)
  - all odd degrees  
(Reztsov, 1990) (Beckers & C., 1993) (C. & Sloan, 1996)
- for  $s = 3$ 
  - degree 5 (Frolov, 1977)



M. Beckers



M.B. Носков



A. Резцов



I.H. Sloan

All known minimal formulas of trigonometric degree are lattice rules,  
except...

All known minimal formulas of trigonometric degree are lattice rules, **except...**

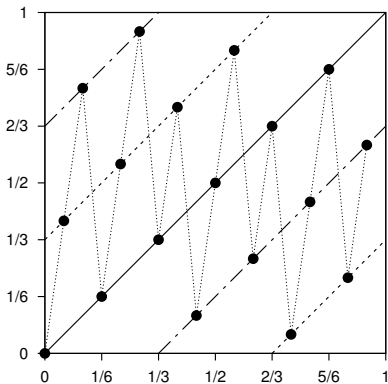
### Theorem (C. & Sloan, 1996)

*The following points*

$$\left( C_p + \frac{j}{2(k+1)}, C_p + \frac{j+2p}{2(k+1)} \right) \begin{array}{l} j = 0, \dots, 2k+1 \\ p = 0, \dots, k \end{array}$$

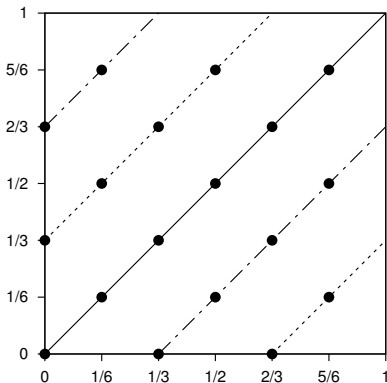
*with  $C_0 = 0$  and  $C_1, \dots, C_k$  arbitrary are the points of a minimal cubature formula of trigonometric degree  $2k+1$ .*





$$k = 2, n = 18, C_1 = \frac{1}{18}, C_2 = \frac{1}{9}$$

$$Q[f] = \frac{1}{n} \sum_{j=0}^{n-1} f\left(\frac{j}{n}, \frac{j(2m+1)}{n}\right) \quad \text{with } n = 2(m+1)^2$$



$k = 2, n = 18, C_1 = C_2 = 0$ : body-centered cubic lattice

$$Q[f] = \frac{1}{2(m+1)^2} \sum_{k=0}^{2m+1} \sum_{j=0}^m f\left(\frac{2j+k}{2(m+1)}, \frac{k}{2(m+1)}\right) \quad \text{with } n = 2(m+1)^2$$

# Technology used to obtain these results:

## Reproducing kernels

The integral  $I$  defines an inner product  $(\phi, \psi) = I[\overline{\phi} \cdot \psi]$ .

Let  $\mathbb{F}$  be a subspace of  $\mathbb{T}^s$ .

Choose  $\phi_1(\mathbf{x}), \phi_2(\mathbf{x}), \dots \in \mathbb{F}$  so that

- $\phi_i(\mathbf{x})$  is  $I$ -orthogonal to  $\phi_j(\mathbf{x})$ ,  $\forall j < i$ , and
- $(\phi_i(\mathbf{x}), \phi_i(\mathbf{x})) = 1$ .

For a given  $k \in \mathbb{N}$  and  $t := \dim(\mathbb{F} \cap \mathbb{T}_k^s)$  we define

$$K(\mathbf{x}, \mathbf{y}) := \sum_{j=1}^t \overline{\phi_j(\mathbf{x})} \cdot \phi_j(\mathbf{y})$$

$K(\mathbf{x}, \mathbf{y})$  is a polynomial in  $2s$  variables of degree  $\leq 2k$ .

## Definition

$K$  is a **reproducing kernel** in the space  $\mathbb{F} \cap \mathbb{T}_k^s$

$$\begin{aligned} \text{if } f \in \mathbb{F} \cap \mathbb{T}_k^s \text{ then } f(\mathbf{a}) &= (f(\mathbf{x}), K(\mathbf{x}, \mathbf{a})) \\ &= \sum_{j=1}^t \phi_j(\mathbf{a}) \cdot I[f(\mathbf{x})\overline{\phi_j(\mathbf{x})}] \end{aligned}$$

The trigonometric monomials form an orthonormal sequence.

$$K(\mathbf{x}, \mathbf{y}) = \sum_{\mathbf{k} \in \Lambda_d} e^{2\pi i \mathbf{k} \cdot (\mathbf{x} - \mathbf{y})}$$

$$\Lambda_d = \left\{ \mathbf{k} \in \mathbb{Z}^s : 0 \leq \sum_{l=1}^s |k_l| \leq \left\lfloor \frac{d}{2} \right\rfloor \right\}$$

A simplifying aspect of the trigonometric case is that the reproducing kernel is a function of one variable:

$$K(\mathbf{x}, \mathbf{y}) = \mathcal{K}(\mathbf{x} - \mathbf{y})$$

with

$$\mathcal{K}(\mathbf{x}') = \sum_{\mathbf{k} \in \Lambda_d} e^{2\pi i \mathbf{k} \cdot \mathbf{x}'}$$

For  $s = 2$  it has the following simple form:

let  $g(z) = \cos(\pi(2\lfloor \frac{d}{2} \rfloor + 1)z) \cos \pi z$ , then

$$\mathcal{K}(\mathbf{x}') = \frac{g(x_1) - g(x_2)}{\sin(\pi(x_1 + x_2)) \sin(\pi(x_1 - x_2))}.$$

# On route to other quality criteria

Assume  $f$  can be expanded into  
an absolutely convergent multiple Fourier series

$$f(\mathbf{x}) = \sum_{\mathbf{h} \in \mathbb{Z}^s} \hat{f}(\mathbf{h}) e^{2\pi i \mathbf{h} \cdot \mathbf{x}}$$

with

$$\hat{f}(\mathbf{h}) = \int_{[0,1]^s} f(\mathbf{x}) e^{-2\pi i \mathbf{h} \cdot \mathbf{x}} d\mathbf{x}$$

Then

$$\begin{aligned} Q[f] - I[f] &= \frac{1}{n} \sum_{j=1}^n \left( \sum_{\mathbf{h} \in \mathbb{Z}^s \setminus \{0\}} \hat{f}(\mathbf{h}) e^{2\pi i \mathbf{h} \cdot \mathbf{x}_j} \right) \\ &= \sum_{\mathbf{h} \in \mathbb{Z}^s \setminus \{0\}} \left( \hat{f}(\mathbf{h}) \frac{1}{n} \sum_{j=1}^n e^{2\pi i \mathbf{h} \cdot \mathbf{x}_j} \right). \end{aligned}$$

Observe that

$$\frac{1}{n} \sum_{j=1}^n e^{2\pi i \mathbf{h} \cdot \mathbf{x}_j} = \begin{cases} 1, & \mathbf{h} \cdot \mathbf{x}_j \in \mathbb{Z} \\ 0, & \mathbf{h} \cdot \mathbf{x}_j \notin \mathbb{Z} \end{cases}$$

A very important tool to investigate the error of a lattice rule is ...

### Definition

The **dual** of the multiple integration lattice  $\Lambda$

$$\Lambda^\perp := \{\mathbf{h} \in \mathbb{Z}^s : \mathbf{h} \cdot \mathbf{x} \in \mathbb{Z}, \forall \mathbf{x} \in \Lambda\}.$$

### Theorem (Sloan & Kachoyan, 1987)

*Let  $\Lambda$  be a multiple integration lattice.*

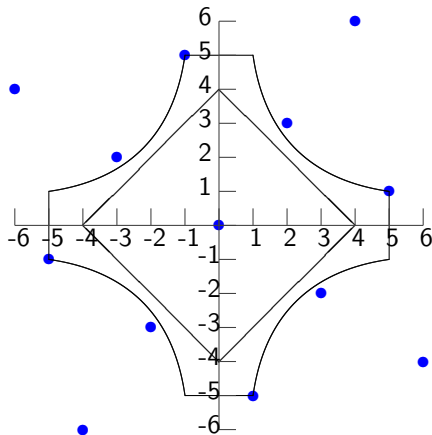
*Then the corresponding lattice rule  $Q$  has an error*

$$Q[f] - I[f] = \sum_{\mathbf{h} \in \Lambda^\perp \setminus \{0\}} \hat{f}(\mathbf{h}).$$



# Example

## Dual lattice of Fibonacci lattice



# Construction criteria

For many years, only used in Russia...

## Definition

The **trigonometric degree** is

$$d(Q) := \min_{\substack{\mathbf{h} \neq \mathbf{0} \\ \mathbf{h} \in \Lambda^\perp}} \left( \sum_{j=1}^s |h_j| \right) - 1.$$

The **enhanced degree**  $\delta := d + 1$ .

Some names:

Mysovskikh (1985–1990), Reztsov (1990), Noskov (1985–1988),  
 Temirgaliev (1991), Semenova (1996–1997), Osipov (2001–2010), Petrov  
 (2004)

# Construction criteria

Mainly used in the 'West'...

## Definition

The **Zaremba index** or **figure of merit** is

$$\rho(Q) := \min_{\substack{\mathbf{h} \neq \mathbf{0} \\ \mathbf{h} \in \Lambda^\perp}} (\bar{h}_1 \bar{h}_2 \cdots \bar{h}_s) .$$

with

$$\bar{h}_j := \begin{cases} 1 & \text{if } h_j = 0 \\ |h_j| & \text{if } h_j \neq 0. \end{cases}$$

Some names:

Maisonneuve (1972), ..., Sloan & Joe (1994), Langtry (1996)

# Where does this come from?

For  $c > 0$  and fixed  $\alpha > 1$ ,  
let  $E_s^\alpha(c)$  be the class of functions  $f$   
whose Fourier coefficients satisfy

$$|\hat{f}(\mathbf{h})| \leq \frac{c}{(\bar{h}_1 \bar{h}_2 \cdots \bar{h}_s)^\alpha},$$

where  $\bar{h} = \max(1, |h|)$ .

Worst possible function in class  $E_s^\alpha(1)$  is

$$f_\alpha := \sum_{\mathbf{h} \in \mathbb{Z}^s} \frac{1}{(\bar{h}_1 \bar{h}_2 \cdots \bar{h}_s)^\alpha} e^{2\pi i \mathbf{h} \cdot \mathbf{x}}$$

$P_\alpha(Q) :=$  the error of the lattice rule for  $f_\alpha$ .

$P_\alpha$  is easy to compute for  $\alpha$  an even integer because  $f_\alpha$  can be written as products of Bernoulli polynomials.

Theoretical convergence is

$$O((\log(n))^{\alpha s} n^{-\alpha}).$$

$P_\alpha$  introduced by (Korobov, 1959)

Obviously related to the figure of merit:

$$\frac{2}{\rho^\alpha} \leq P_\alpha.$$

Figure of merit used by (Maisonneuve, 1972)

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Other criteria:

- $R(\mathbf{z}, n)$  (Niederreiter, 1987)

$$P_\alpha(\mathbf{z}, n) < R(\mathbf{z}, n)^\alpha + \mathcal{O}(n^{-\alpha})$$

- Discrepancy

$$D_N = O\left(\frac{(\log N)^{s-1}}{\rho}\right)$$



H. Niederreiter

# Yet another way to look at this

Assume  $f$  can be expanded into  
an absolutely convergent multiple Fourier series

$$f(\mathbf{x}) = \sum_{\mathbf{h} \in \mathbb{Z}^s} \hat{f}(\mathbf{h}) e^{2\pi i \mathbf{h} \cdot \mathbf{x}} \quad \text{with} \quad \hat{f}(\mathbf{h}) = \int_{[0,1]^s} f(\mathbf{x}) e^{-2\pi i \mathbf{h} \cdot \mathbf{x}} d\mathbf{x}$$

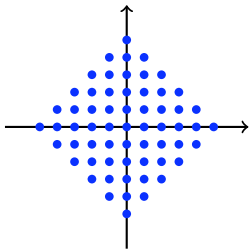
- Mark **region of interest**  $\mathcal{A}_s(m)$  in Fourier domain of “degree”  $m$ .
- Ask to integrate those Fourier terms exactly, i.e.

$$\Lambda^\perp \cap \mathcal{A}_s(m) = \{\mathbf{0}\}.$$

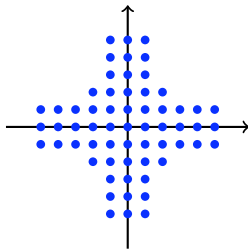
- $\Rightarrow$  Rule of degree (at least)  $m$ .
- Different regions  $\mathcal{A}_s(m)$  possible:
  - Trigonometric degree.
  - Zaremba cross degree.
  - Product trigonometric degree.
  - ...

# Corresponding Fourier spectra

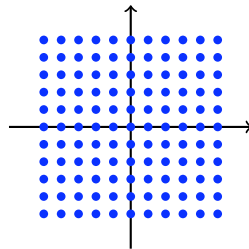
Take  $m = 5$  (and  $s = 2$ ):



Trigonometric degree



Zaremba degree



Product degree

For  $s \rightarrow \infty$  these shapes grow exponentially.  
Consequently the number of nodes grows exponentially.



Modern interpretation of  $P_\alpha$  is the squared **worst-case error** in a RKHS with Korobov kernel with smoothness  $\alpha$ .

In general, for a shift-invariant kernel  $K$  and rank-1 lattice points

$$e^2(\Lambda, K) = - \int_{[0,1]^s} K(\mathbf{x}, \mathbf{0}) \, d\mathbf{x} + \frac{1}{n} \sum_{k=0}^{n-1} K\left(\left\{\frac{k\mathbf{z}}{n}\right\}, \mathbf{0}\right)$$

see e.g. (Hickernell, 1998)

Typical form for a weighted space:

$$e_s^2(\mathbf{z}) = -1 + \frac{1}{n} \sum_{k=0}^{n-1} \prod_{j=1}^s \left[ 1 + \gamma_j \omega\left(\left\{\frac{kz_j}{n}\right\}\right) \right]$$



This is a tensor product space: a **product** of 1-dimensional kernels

The **weights**  $\gamma_j$ ,  $\gamma_1 \geq \gamma_2 \geq \dots \geq \gamma_s$ , model anisotropicness of the integrand functions

Between the **big braces** we have the 1-dimensional kernel

# Searches for lattice rules

Remember that

- 1 The cost to verify that a lattice rule has degree  $d$  is proportional to  $d^s$ , so only “moderate” dimensions are feasible.
- 2 The search space is huge.

⇒ Restrict the search space.

For example:

### Definition

A rank-1 simple lattice is generated by one vector  $\mathbf{z}$  and has the form

$$Q[f] := \frac{1}{n} \sum_{j=0}^{n-1} f \left( \left\{ \frac{j\mathbf{z}}{n} \right\} \right)$$

$$P_n := \left\{ \left\{ \frac{j\mathbf{z}}{n} \right\} : j = 0, \dots, n-1 \right\}, \quad \mathbf{z} \in U_n^s.$$

# For example:

## Definition

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$$Q[f] := \frac{1}{n} \sum_{j=0}^{n-1} f \left( \left\{ \frac{j\mathbf{z}}{n} \right\} \right)$$

- Restricting to rank-1 simple rules  
→ only 1 vector,  $s - 1$  components, to be determined.
- Further restriction of the search space:  
consider only generator vectors of the form

$$\mathbf{z}(\ell) = (1, \ell, \ell^2 \pmod n, \dots, \ell^{s-1} \pmod n), 1 \leq \ell < n$$

(Korobov, 1959)

## Technology used: matrices

Any  $s$ -dimensional lattice  $\Lambda$  can be specified in terms of  $s$  linearly independent vectors  $\{\mathbf{a}_1, \mathbf{a}_2, \dots, \mathbf{a}_s\}$ .

→ These vectors are known as **generators** of  $\Lambda$ .

Associated with the generators is an  $s \times s$  **generator matrix**  $A$  whose rows are  $\mathbf{a}_1, \mathbf{a}_2, \dots, \mathbf{a}_s$ .

All  $\mathbf{h} \in \Lambda$  are of the form  $\mathbf{h} = \sum_{i=1}^s \lambda_i \mathbf{a}_i = \boldsymbol{\lambda} A$  for some  $\boldsymbol{\lambda} \in \mathbb{Z}^s$ .

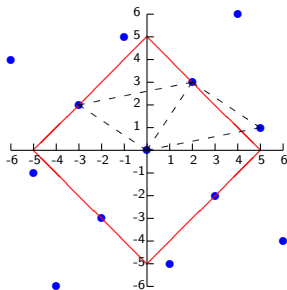
The dual lattice  $\Lambda^\perp$  may be defined as having generator matrix  $B = (A^{-1})^T$ .

It can be shown that the number of points  $n = |\det A|^{-1} = |\det B|$ .

# Recent searches for low dimensions: *K*-Optimal rules

*Not* restricted to rank-1 lattices.

Based on a property of the dual lattice:



Argument by (C. & Lyness, 2001):

*It is reasonable to believe that the lattice  $\Lambda$  of an optimal lattice rule will have  $\Lambda^\perp$  with many elements on the boundary of  $\text{conv}S(O_s, d + 1)$  (a scaled version of the unit octahedron).*

High computational cost,  $\mathcal{O}(\delta^{s^2-1})$ .  $(\delta = d + 1)$

- (C. & Lyness, *Math. Comp.*, 2001): 3D ( $\delta \leq 30$ , 4D ( $\delta \leq 24$ ))
- (Lyness & Sørøvik, *Math. Comp.*, 2006): 5D ( $\delta \leq 12$ )

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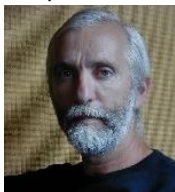
Restricting the search to (skew-)circulant generator matrices, reduces the cost to  $\mathcal{O}(\delta^{2s-2})$ .

- (Lyness & Sørøvik, *Math. Comp.*, 2004): 4D
- (C. & Govaert, *J. Complexity*, 2003): 5D, 6D

This also lead to closed expressions for arbitrary degrees.



J. Lyness



T. Sørøvik



H. Govaert



# Packing factor

## Definition

The **packing factor**

$$\hat{\rho}(n) := \frac{\delta^s}{s!n}.$$

This is a measure of the efficiency of a rule.

It is convenient for making pictures because  $0 \leq \hat{\rho}(n) \leq 1$ .

Actually,  $\hat{\rho}(n)$  is bounded above by  
 the **density of the densest lattice packing**  
 of the crosspolytope (octahedron)  $\theta(O_s)$ .  
 (→ link with “Geometry of numbers”)

Known values:

- $\theta(O_1) = \theta(O_2) = 1$
- $\theta(O_3) = \frac{18}{19}$  (Minkowski, 1911) used by (Frolov, 1977)

This provides a (higher) lower bound for lattice rules for trigonometric degree:

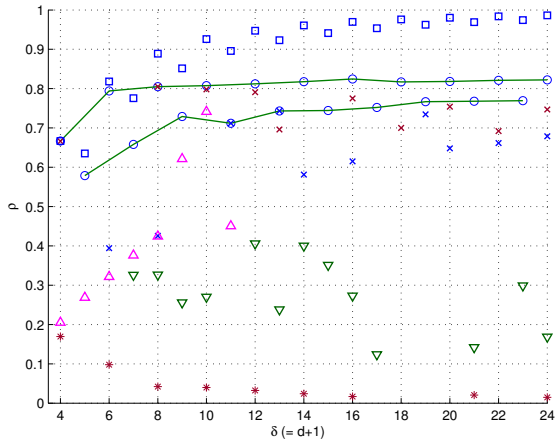
$$n \geq \frac{(d+1)^s}{s! \theta(O_s)}.$$

Lattice rules provide constructive lower bounds for  $\theta(O_s)$ . From a lattice rule with  $n$  points follows

$$\theta(O_s) \geq \frac{(d+1)^s}{s! n}.$$

The best known bounds for  $\theta(O_4)$ ,  $\theta(O_5)$  and  $\theta(O_6)$  come from lattice rules (C., *East Journal on Approximations*, 2006).

## Results: 4D



- refers to  $n_{KO}$ , □ refers to  $n_{ME}$ .
- × refers to (Noskov & Semenova, 1996)+corrections
- \* refers to (C., Novak & Ritter, 1999)
- × refers to (Temirgaliev, 1991), △ refers to Good lattices
- ▽ refers to Korobov rules (Maisonneuve, 1972)

# K-optimal rules: conclusions

- The search for  $K$ -optimal lattice rules is expensive.
- The **packing factor** is related to the concept **critical lattice** (a global minimum)  
As a side effect it delivered the best known constructive lower bounds for  $\theta(s)$ , for  $s = 4, 5, 6$ .
- There are also local minima for the determinant of admissible lattices  
→ **extremal lattices**  
The corresponding lattices can be used to bootstrap the construction of higher degree lattice rules (in no-time) and sequences.
- More recent: approach based on Golomb rules (Sørøvik, MCQMC2012)

# Recent searches for higher dimensions: Component-by-component construction

- Focus on **rank-1 lattice rules**  $\Rightarrow$  find 1 vector  $\mathbf{z}$ .
- Idea: search  $\mathbf{z}$  component by component

2000: I. Sloan & A. Reztsov (Tech. Report)  
published Math. Comp. 2002  
unweighted Korobov space,  $n$  prime

Note that Korobov (1959) presented a constructive proof using the CBC-principle.



I.H.Sloan



A. Reztsov

# Some milestones of component-by-component

- 2000-2002: F. Kuo (PhD) with S. Joe  
weighted Korobov space, weighted Sobolev space
- MCQMC 2002: J. Dick & F. Kuo  
basically for weighted Korobov space,  $n$  a product of few primes, but  
partial search, faster and for millions of points
- MCQMC 2004, 2006: D. Nuyens & C.  
fast construction in  $O(sn \log(n))$ , basic case for  $n$  prime, but also  
possible for any composite  $n$  (and full search)



F. Kuo



S. Joe



J. Dick



D. Nuyens

# The CBC algorithm in a shift-invariant RKHS

**for**  $s = 1$  **to**  $s_{\max}$  **do**  
     **for all**  $z$  **in**  $U_n$  **do**

$$e_s^2(z) = -1 + \frac{1}{n} \sum_{k=0}^{n-1} \prod_{j=1}^s \left[ 1 + \gamma_j \omega \left( \left\{ \frac{kz_j}{n} \right\} \right) \right]$$

**end for**

$$z_s = \operatorname{argmin}_{z \in U_n} e_s^2(z)$$

**end for**

Computational cost:  $O(s_{\max} n^2)$

# Rephrasing CBC: matrix-vector form

The inner loop can be formulated as a matrix-vector product with matrix

$$\Omega_n := \left[ \omega \left( \left\{ \frac{kz}{n} \right\} \right) \right]_{\substack{z \in U_n \\ k \in \mathbb{Z}_n}} = \left[ \omega \left( \frac{k \cdot z \bmod n}{n} \right) \right]_{\substack{z \in U_n \\ k \in \mathbb{Z}_n}}$$

This matrix has a lot of structure!

A matrix-vector multiplication can be done in  $O(n \log n)$   
(Nuyens & C. 2005, 2006)

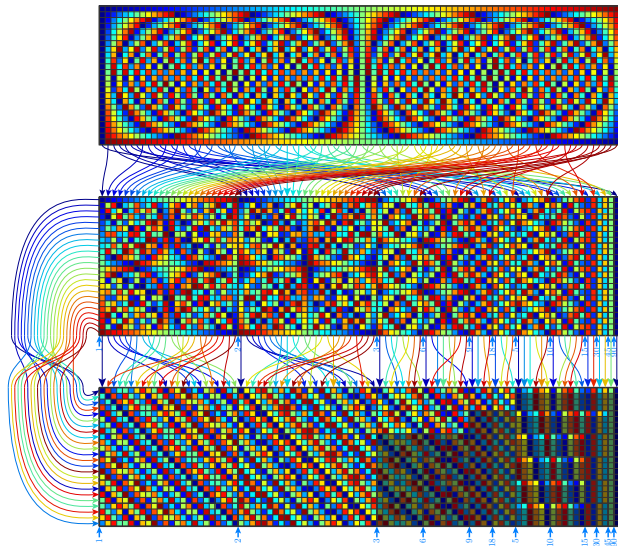
⇒ Construction then takes  $O(sn \log n)$  using  $O(n)$  memory

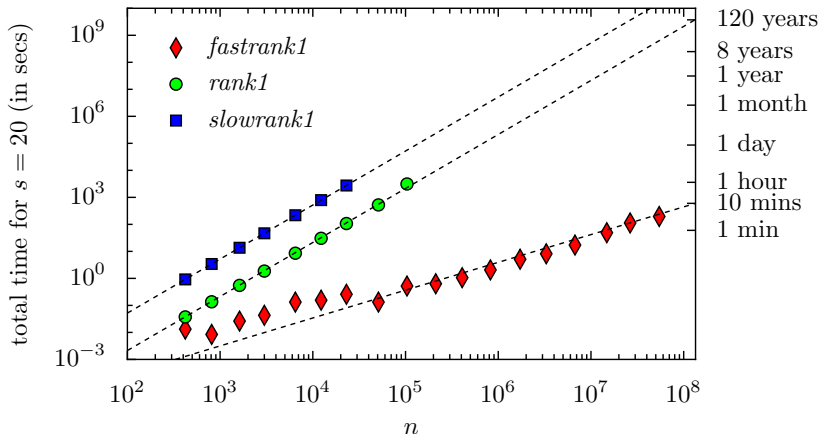


# An example matrix $\Omega_n$ and its permutations

A nice view on  
 $90 = 2 \times 3^2 \times 5$

The blocks of  
the last matrix  
are  
diagonalizable  
by FFT's



Results in  $O(sn \log(n))$ 

Timings anno 2004 for 20 dimensions  
generated on a P4 2.4GHz ht, 2GB RAM

# Combination of approaches

Inspired by “classical” approach and

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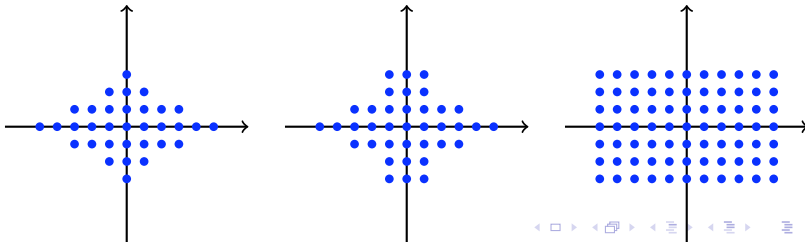
H. Woźniakowski

I.H. Sloan

weighted spaces from QMC (Sloan & Woźniakowski, 1998),

→ “weighted degree of exactness”:

For example:



# A new worst case setting

Amend the Korobov space  $E_\alpha$  to make new space  $H$  with reproducing kernel

$$K(\mathbf{x}, \mathbf{y}) = \sum_{\mathbf{h} \in \mathcal{A}_s(m)} \exp(2\pi i \mathbf{h} \cdot (\mathbf{x} - \mathbf{y})) + \sum_{\mathbf{h} \notin \mathcal{A}_s(m)} \frac{\exp 2\pi i \mathbf{h} \cdot (\mathbf{x} - \mathbf{y})}{r_\alpha(\boldsymbol{\gamma}, \mathbf{h})}.$$

The squared worst case error of a rank-1 lattice rule is now

$$e_{n,s}^2(\mathbf{z}) = \sum_{\substack{\mathbf{0} \neq \mathbf{h} \in \mathcal{A}_s(m) \\ \mathbf{h} \cdot \mathbf{z} \equiv 0 \pmod{n}}} 1 + \sum_{\substack{\mathbf{h} \notin \mathcal{A}_s(m) \\ \mathbf{h} \cdot \mathbf{z} \equiv 0 \pmod{n}}} \frac{1}{r_\alpha(\boldsymbol{\gamma}, \mathbf{h})}.$$

→ CBC-algorithm (C., Kuo & Nuyens, 2010)

# Error estimation

In practice one wants more than 1 approximation.

Common approaches (for all types of cubature):

- randomization (randomly shifted rules) (Cranley & Patterson, 1976)

# Error estimation

In practice one wants more than 1 approximation.

Common approaches (for all types of cubature):

- randomization (randomly shifted rules) (Cranley & Patterson, 1976)
- embedded sequences
  - copy rules, with intermediate lattice rules (Joe & Sloan, 1992)  
augmentation sequences (Li, Hill & Robinson, 2007)
  - embedded rank-1 rules  
(Hickernell, Hong, L'Ecuyer, Lemieux, *SISC* 2000)  
(C., Kuo, Nuyens, *SISC* 2006)  
(C. & Nuyens, *MCQMC2008*)



T. Patterson



R. Hong

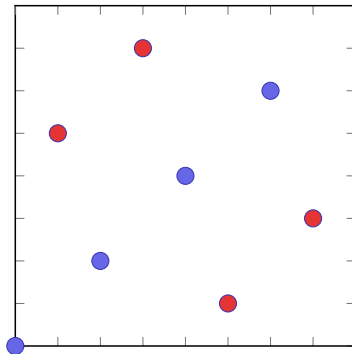


P. L'Ecuyer



C. Lemieux

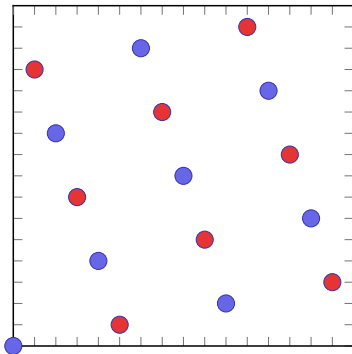
# Example of embedded rank-1 rules



$n = 8$

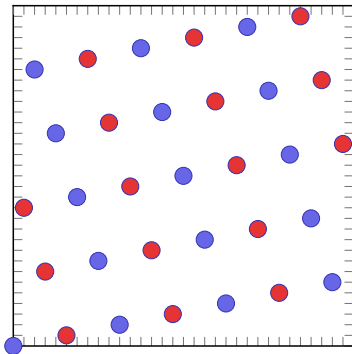


# Example of embedded rank-1 rules



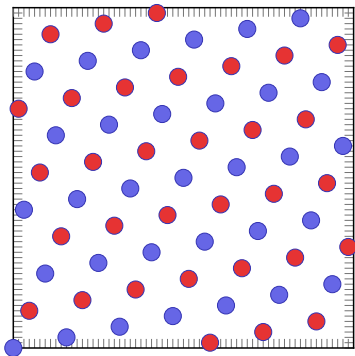
$n = 16$

# Example of embedded rank-1 rules



$n = 32$

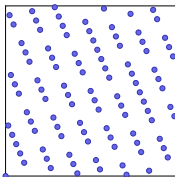
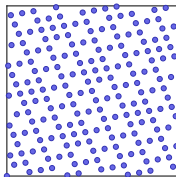
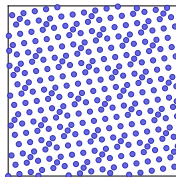
# Example of embedded rank-1 rules



$n = 64$

# This is not restricted to powers of 2

- The structure of the points using Gray code or radical inverse ordering is similar to that of a net. The unit cube gets filled with smaller lattices which consists of smaller lattices and so on.
- Starting from a good lattice sequence we can stop anywhere and have a good uniform distribution (Hickernell, Kritzer, Kuo, Nuyens, 2011)

 $n = 100$  $n = 200$  $n = 300$ 

P. Kritzer

# Is the Weyl sequence a relative?

- Simple rank-1 lattice:

$$\mathbf{x}^{(k)} = \left\{ \frac{k \mathbf{z}}{n} \right\}, \text{ for } k = 0, 1, 2, \dots, n - 1.$$

- Embedded rank-1 lattice: in order to stop at any time, you need a good ordering of the points:

$$\mathbf{x}^{(k)} = \left\{ \frac{\varphi(k)}{n} \mathbf{z} \right\}, \text{ for } k = 0, 1, 2, \dots, n - 1.$$

- If  $n$  is very large, this can be seen as an extensible cubature rule.
- Weyl sequence: Take  $n \rightsquigarrow \infty$ , then  $\ell/n$  has an infinite digit expansion, i.e. think “irrational”.

Now group on  $\mathbf{z}/n$ , and take each  $z_j/n = \xi_j$  an irrational:

$$\mathbf{x}^{(k)} = \{k \boldsymbol{\xi}\}, \text{ for } k = 0, 1, 2, \dots$$

This could be interpreted as an infinite extensible “lattice”.

# Weyl sequence for periodic functions

Introduce weights and achieve higher order of convergence for periodic functions.

(Niederreiter, 1973) (Sugihara & Murota, 1982)

(Vandewoestyne, C. & Warnock, 2007)

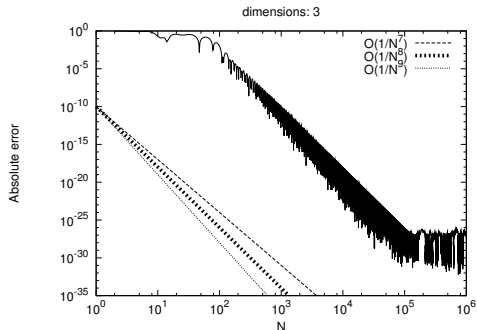


M. Sugihara



B. Vandewoestyne

Example: 3D,  $\mathcal{O}(n^{-8})$



# Final remarks

## Construction:

- Searches for lattice rules using the “classical” criteria are doomed to fail for increasing dimensions.
- The CBC algorithm, relying on “worst-case-error” for “reproducing kernel Hilbert spaces” beats this curse of dimensionality. Rules can be constructed very fast even if  $n$  and  $s$  are large.

But work remains to be done, e.g.,

- for CBC, tuning of the function space using the weights,
- practical error estimates based on sequences.

Finally note that

- lattice rules are useful for low and high dimensions, and are not only for integrating periodic functions;
- all quality criteria have a reason to exist;
- the difference between lattice rules and “classical” low discrepancy sequences evaporates.

Lattice rules with large  $n$  can be constructed easily and can be used as sequences.



Finally note that

- lattice rules are useful for low and high dimensions, and are not only for integrating periodic functions;
- all quality criteria have a reason to exist;
- the difference between lattice rules and “classical” low discrepancy sequences evaporates.

Lattice rules with large  $n$  can be constructed easily and can be used as sequences.

Use a lattice rule anywhere & anytime!

This was a story about integration but the above suggestion also applies to you if you are involved in approximation.

The end.  
Thank you!

# The end.

Thank you!

A special “thank you” to those that put their picture on the web.  
Don't forget to update it!